

Use of Scilab to demonstrate concepts in linear algebra and polynomials

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- 1 Matrices
- 2 Polynomials
- 3 Coprime polynomials
- 4 Fourier transform, polynomials, matrices

- Scilab is free.
- Matrix/loops syntax is same as for Matlab.
- Scilab provides all basic and many advanced tools.
- This talk focus: linear algebra and polynomials.

Defining a matrix

- `A=[1 3 4 6]`
- `B=[1 3 4 6;5 6 7 8]`
- `size(A)`, `length(A)`, `ones(A)`, `zeros(B)`, `zeros(3,5)`

- `A=rand(3,3)`
- `det(A)`, `spec(A)`, `trace(A)`
- `sum(spec(A))`
- ```
if sum(spec(A))==trace(A) then
 disp('yes, trace equals sum')
else
 disp('no, trace is not sum ')
end
```
- `prod(spec(A))-det(A)`

## (Block) diagonalize A?

Let  $A$  be a square matrix ( $n \times n$ ) with distinct eigenvalues  $\lambda_1, \dots, \lambda_n$ . Eigenvectors (column vectors)  $v_1$  to  $v_n$  are then independent.

$$Av_1 = \lambda_1 v_1 \quad Av_2 = \lambda_2 v_2 \quad \dots \quad Av_n = \lambda_n v_n$$

$$A [v_1 \ v_2 \ \dots \ v_n] = [v_1 \ v_2 \ \dots \ v_n] \begin{bmatrix} \lambda_1 & 0 & \dots & 0 \\ 0 & \lambda_2 & \dots & 0 \\ \vdots & \vdots & \ddots & \vdots \\ 0 & 0 & \dots & \lambda_n \end{bmatrix}$$

(Column scaling of vectors  $v_1$ , etc is just post-multiplication.)

- `[spe, vect] = spec(A)`
- `inv(vect)*A*vect`

Inverse exists because of independence assumption on eigenvectors. Use 'bdiag' command for block diagonalization (when non diagonalizable).

- $\text{rank}(A) = \text{svd}(A)$
- $[u, s, v] = \text{svd}(A)$
- check  $u' \cdot \text{inv}(u) \quad u*s*v - A$

## Example: Income tax

Income tax for a man earning Rs. NET (after exempted deductions) is

0% for the first 1,50,000

10% for the part between 1,50,000 and 3,00,000

20% for the part between 3,00,000 and 5,00,000

30% for the part above 5,00,000

- 
- $[u, s, v] = \text{svd}(A)$
- check  $u' \text{-inv}(u)$        $u*s*v-A$



# Defining polynomials

Polynomials play a very central role in control theory: transfer functions are ratio of polynomials.

- `s=poly(0,'s')`      `s=poly(0,'s','roots')`
- `p=s^2+3*s+2`      `p=poly([2 3 1],'s','coeff')`
- `roots(p)`      `horner(p,5)`
- `a = [1 2 3]`      `horner(p,a)`      `horner(p,a')`
- `w=poly(0,'w')`      `horner(p,%i*w)`

- `p=poly([1 2 3 4 -3], 's', 'coeff')`
- `cfp=coeff(p)`
- `diffpcoff=cfp(2:length(cfp)).*[1:length(cfp)-1]`
- `diffp=poly(diffpcoff, 's', 'coeff')`
- `degree(p)` can be used instead of `length(cfp)-1`

## More about horner

- `w=poly(0,'w')`      `horner(p,(1+w)/(1-w))`
- `a=-rand(1,4); p=poly(a,'s');`
- `q=horner(p,(w-1)/(1+w))` // bilinear(Cayley transform)
- `abs(roots(numer(q)))`

# Multiplication and convolution

Output of a (linear and time-invariant) dynamical system is the convolution of the input signal with the 'impulse response'.

Convolution: central role.

Polynomial multiplication is related to convolution of their coefficients

- `a=[1 2 3]; b=[4 5 6]; convol(a,b)`
- `pa=poly(a,'s','coeff'); pb=poly(b,'s','coeff');`  
`coeff(pa*pb)`

To convolve  $u(\cdot)$  by  $h(\cdot)$  is a linear operation on  $u(\cdot)$ .

Write  $h(s) = h_0 + h_1s + h_2s^2 + \dots + h_ns^n$  (similarly  $u(s)$ )

convolution  $y := h * u$  (convolution of  $h$  and  $u$ ).

$$y(k) = \sum_{j=0}^{n+m} h(j)u(k-j) \quad (u \text{ has degree } m).$$

# Matrix for convolution

$$[y_0 \ y_1 \ \cdots \ y_{n+m}] = [u_0 \ u_1 \ \cdots \ u_m] C_h$$

where the matrix  $C_h$  with  $m + 1$  rows and  $n + m + 1$  columns is defined as

$$\begin{bmatrix} h_0 & h_1 & h_2 & \cdots & h_n & 0 & \cdots & 0 \\ 0 & h_0 & h_1 & \cdots & h_{n-1} & h_n & \cdots & 0 \\ \vdots & \ddots & \ddots & \cdots & \ddots & \ddots & \cdots & \vdots \\ 0 & 0 & 0 & \cdots & & & & h_n \end{bmatrix}$$

# Coprime polynomials

Numerator and denominator polynomials of a transfer function being coprime is critical for controllability and observability of dynamical systems: Kalman

- polynomials  $a(s)$  and  $b(s)$  are called coprime if they have no common root.
- Equivalently, their gcd (greatest common divisor) is 1.

Problem: given  $a(s)$  and  $b(s)$ , find polynomials  $p(s)$  and  $q(s)$  such that  $ap + bq = 1$ .

Easy: take  $p := -b$  and  $q = a$ . Relation to coprimeness??

## Coprimeness: equivalent statements

Consider polynomials  $a(s)$  (of degree  $m$ ) and  $b(s)$  (of degree  $n$ ).  
Following statements are equivalent.

- $a$  and  $b$  are coprime.
- there exist no polynomials  $p$  and  $q$  with  $\text{degree}(p) < n$  and  $\text{degree}(q) < m$  such that  $ap + bq = 0$ .
- there exist polynomials  $u$  and  $v$  such that  $au + bv = 1$  (their gcd).

In fact,  $u$  and  $v$  having degrees at most  $n - 1$  and  $m - 1$  respectively can be found. Then they are unique.

# Equivalent matrix formulations

$ap + bq$  can be considered as having coefficients obtained from

$$[p_0 \ p_1 \ \cdots \ p_{m-1} \ q_0 \ q_1 \ \cdots \ q_{n-1}] \begin{bmatrix} C_a \\ C_b \end{bmatrix}$$

$C_a$  and  $C_b$  have  $m$  rows and  $n$  rows respectively, and both have  $m + n$  columns each.

Hence, following are equivalent

- $a$  and  $b$  are coprime
- $\begin{bmatrix} C_a \\ C_b \end{bmatrix}$  is nonsingular
- $\begin{bmatrix} C_a \\ C_b \end{bmatrix}$  has  $[1 \ 0 \ \cdots \ 0]$  in its (left)-image

Above matrix: Sylvester resultant matrix, its determinant:  
resultant of two polynomials



- `sylvester_mat.sci` function constructs required matrix
- `linsolve(sy',[1 0 0 ... 0]')` // for coprime and non-coprime
- `[x,kern]=linsolve(sy',[0 0 0 ... 0]')`

Uncontrollable and unobservable modes are related to eigenvectors corresponding to the eigenvalue which is the 'common root'.

## Command 'find' : extracts TRUE indices

Consider the following problem: polynomials  $a$  and  $b$  might have roots 'close by'.

(Very difficult to control/observe: very high energy or input levels needed to control, or measurement very sensitive to noise. This is due to 'close to' uncontrollable/unobservable).

Find which are close to each other.

- Find roots of  $a$  and  $b$ . For each root of  $a$ , check if a root of  $b$  is within specified tolerance 'toler'.
- Two for loops?
- find allows extraction of indices satisfying a boolean expression

## Coefficients and powers as vectors

Evaluation of a polynomial  $p(s)$  at a value  $s = a$  is a linear map on the coefficients.

$$p(a) = [p_0 \ p_1 \ \cdots \ p_n][1 \ a \ a^2 \ \cdots \ a^n]'$$

Moreover, if  $n$  and  $m$  are the degrees of  $p$  and  $q$  respectively,

$$p(s)q(s) = [1 \ s \ s^2 \ \cdots \ s^n] \begin{bmatrix} p_0 \\ p_1 \\ \vdots \\ p_n \end{bmatrix} [q_0 \ q_1 \ \cdots \ q_m] \begin{bmatrix} 1 \\ s \\ \vdots \\ s^m \end{bmatrix}$$

$(n + 1) \times (m + 1)$  matrix.

# Bezoutian matrix

Bezoutian of a pair of polynomials  $p(s)$  and  $q(s)$  is defined as the symmetric matrix  $B$  such that (suppose  $n \geq m$ )

$$[1 \ x \ x^2 \ \dots \ x^{n-1}] B \begin{bmatrix} 1 \\ y \\ \vdots \\ y^{n-1} \end{bmatrix} = \frac{p(x)q(y) - p(y)q(x)}{x - y}$$

$B$  is  $n \times n$  matrix.  $B$  is nonsingular if and only if  $p$  and  $q$  are coprime.

Size of  $B$  is roughly half the size of the Sylvester resultant matrix.

$B$  is symmetric.

(Both can have polynomials (in  $\gamma$ ) as their coefficients.)

# Discrete Fourier Transform

For a periodic sequence: DFT (Discrete Fourier Transform) gives the frequency content.

Linear transformation on the input sequence.

Take signal values of just one period: finite dimensional signal (due to periodicity of  $N$ ).

$$X(k) := \sum_{n=0}^{N-1} x(n)e^{\frac{-2\pi ik}{N}n} \text{ for } k = 0, \dots, N-1 \text{ (analysis equation)}$$

$e^{-2\pi ik/N}$  is the  $N^{\text{th}}$  root of unity.

Inverse DFT for the synthesis equation. Normalization constants vary in the literature.

# Discrete Fourier Transform

What is the matrix defining relating the DFT  $X(k)$  of the signal  $x(n)$ ? Define  $\omega := e^{\frac{-2\pi ik}{N}n}$ .

$$\begin{bmatrix} X(0) \\ X(1) \\ \vdots \\ X(N-1) \end{bmatrix} = \begin{bmatrix} 1 & 1 & 1 & \cdots & 1 \\ 1 & \omega & \omega^2 & \cdots & \omega^{N-1} \\ 1 & \omega^2 & \omega^4 & \cdots & \omega^{2N-2} \\ \vdots & \vdots & \vdots & \ddots & \vdots \\ 1 & \omega^{N-1} & \omega^{2N-2} & \cdots & \omega^{(N-1)^2} \end{bmatrix} \begin{bmatrix} x(0) \\ x(1) \\ \vdots \\ x(N-1) \end{bmatrix}$$

(Note:  $\omega^N = 1$ , etc.)

Check that the above  $N \times N$  matrix has nonzero determinant.

(Change of basis.) Moreover, columns are orthogonal.

Orthonormal? (Normalization (by  $\sqrt{N}$ ) not done yet.)

# Discrete Fourier Transform and interpolation

Van der monde matrix: closely related to interpolation problems

Of course, inverse DFT is nothing but interpolation! Used in computation of determinant of a polynomial matrix.

Construct  $p(s) := x_0 + x_1s + x_2s^2 \cdots + x_{N-1}s^{N-1}$

To obtain  $X(k)$ , evaluate  $p$  at  $s = \omega^k$ .

$X(k) = p(\omega^k)$  horner command

Given values of  $p(\omega^k)$  for various  $\omega^k$  (i.e.,  $X(k)$ ), find the coefficients of the polynomial  $p(s)$ : inverse DFT: interpolation of a polynomial to 'fit' given values at specified (complex) numbers.

Since many powers of  $\omega$  are repeated in that matrix (only  $N - 1$  powers are different, many real/imaginary parts are repeated for even  $N$ ), redundancy can be **drastically** decreased.  
Length of the signal is a power of 2: recursive algorithm possible.



# FFT: recursive implementation

- Separate  $p(s)$  (coefficients  $x_0, \dots, x_{N-1}$ ) into its even and odd powers (even and odd indices  $k$ ).  $N$  is divisible by 2.
- Compute DFT of  $p_{\text{odd}}$  and  $p_{\text{even}}$  separately. (Do same separation, if possible.)
- Let  $X_{\text{odd}}$  and  $X_{\text{even}}$  denote the individual DFT's. (Same length.)
- Define  $D := \text{diag}(1, \omega, \omega^2, \dots, \omega^{\frac{N}{2}-1})$
- Combine the two separate DFT's using the formula

$$\begin{aligned} X(k) &= X_{\text{even}} + DX_{\text{odd}} \text{ for } k = 0, \dots, \frac{N}{2} - 1 \\ X(k) &= X_{\text{even}} - DX_{\text{odd}} \text{ for } k = \frac{N}{2}, \dots, N - 1 \end{aligned}$$

- Matrices and polynomials provide rich source of problems
- With good computational tools, the future lies in computational techniques
- Scilab provides handy tools
- We saw: `if elseif else end for horner poly coeff`
- Recursive use of function
- `find conv max min`